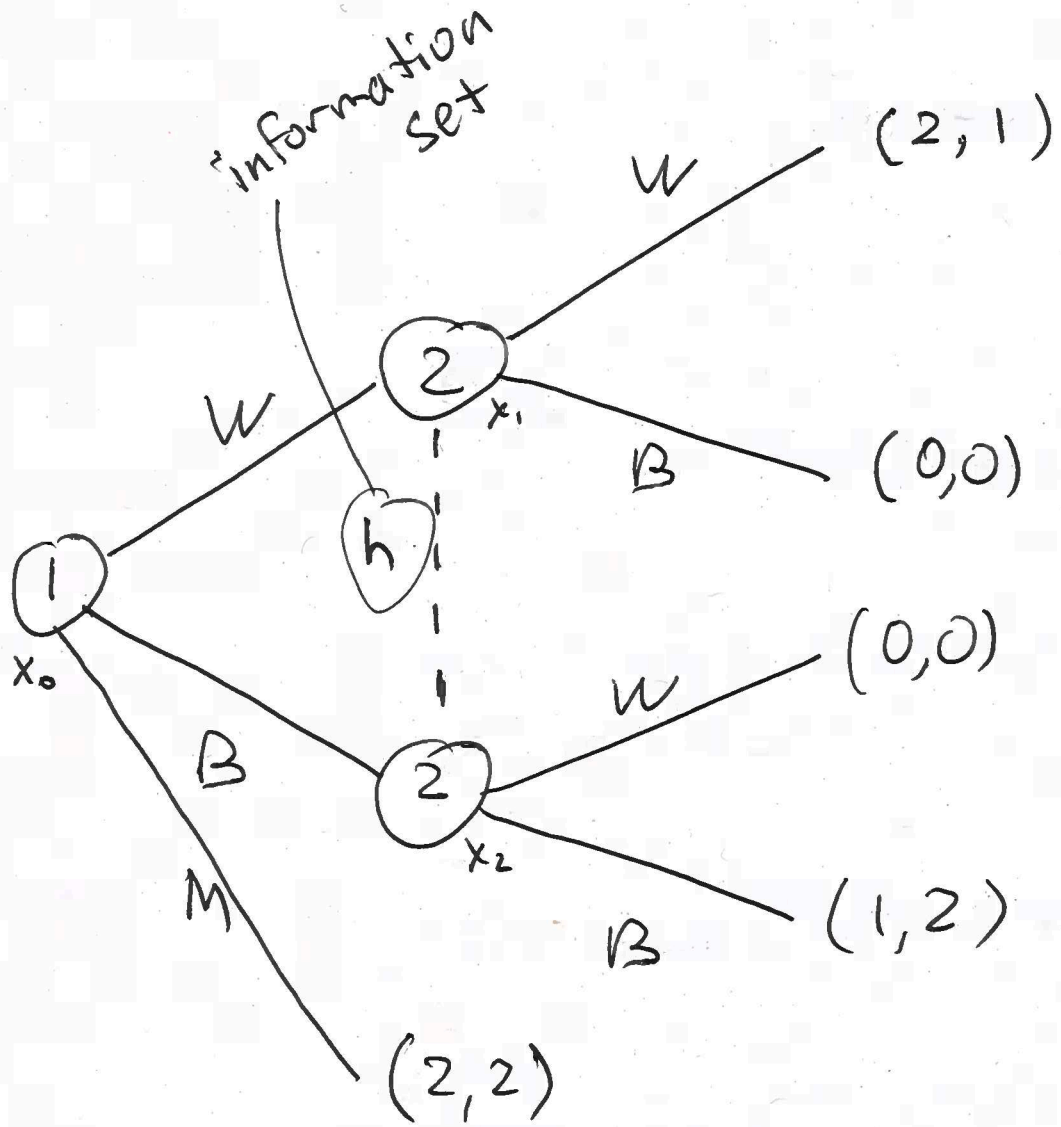


Extensive - Form Games

Games in Tree Form

↗ sequential
↘ simultaneous

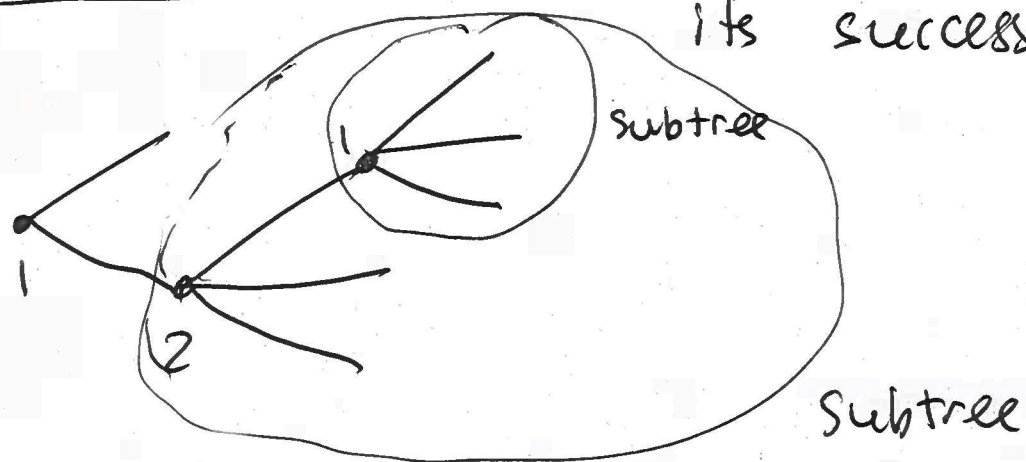


Ingredients of an Extensive-Form Game

- ① Players $0, 1, 2, \dots, N$ Player 0 is Nature
- ② Nodes
- Decision Nodes ($x \in X$, initial node x_0)
 - Terminal Nodes ($z \in Z$)

- ③ The tree:
- each node has a predecessor (except for x_0)
 - each decision node has two or more successors
 - each terminal node has no successors.

Subtree (or branch) is any decision node and all its successors and all successors of successors, etc.



(4) Decision mapping $i: X \rightarrow \{0, 1, \dots, N\}$

that specifies which player gets to move
in each decision node

$i(x) \equiv$ player who
moves in node x

(5) Actions: $A(x)$ is the set of actions
available to player $i(x)$ at decision node $x \in X$

(6) Information sets: What a player knows about
moves made earlier in the game.

Info sets: Let M be a partition of the
decision nodes

↳ A player cannot distinguish between the decision
nodes within an information set she is at,
but can distinguish between info. sets

$h(x)$ = info. set that contains node x

Constraints:

(a) In an info. set, the same player moves in each node of that set

(b) In an info. set, the same actions are available in each node of that set.

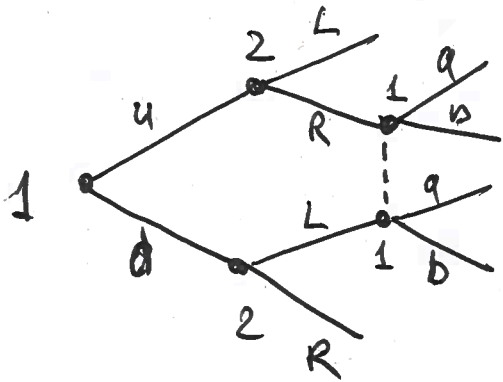
if $x' \in h(x)$, then $i(x') = i(x)$ } write $i(h)$
if $x' \in h(x)$, then $A(x') = A(x)$ } write $A(h)$

⑦ Nature's moves: Non-strategic, random according to a commonly known probability distribution

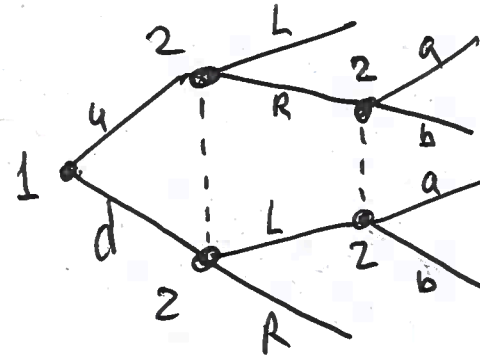
⑧ Payoffs: For each player i , $u_i: Z \rightarrow \mathbb{R}$ specifies player i 's payoff $u_i(z)$ at each terminal node $z \in Z$.

Terminology

1. A game of perfect recall: Players don't forget anything they have observed or known.



Player 1 has imperfect recall



Player 2 has imperfect recall

2. A game of perfect information = "no hidden moves". Each info. set contains only one node. Players know the entire history of previous moves.

→ Simultaneous-move game is a game of imperfect information

3. A game of complete information:

The structure of the game (in particular, the players' payoffs) is common knowledge.

Strategies

Let H_i be the collection of info. sets of player i :

$$H_i = \{h : i(h) = i\}$$

Let A_i be the set of actions of player i in the game tree:

$$A_i = \bigcup_{h \in H_i} A(h)$$

A pure strategy of player i is a function

$$s_i : H_i \rightarrow A_i \quad \text{with the property that} \\ s_i(h) \in A(h) \quad \text{for all } h \in H_i$$

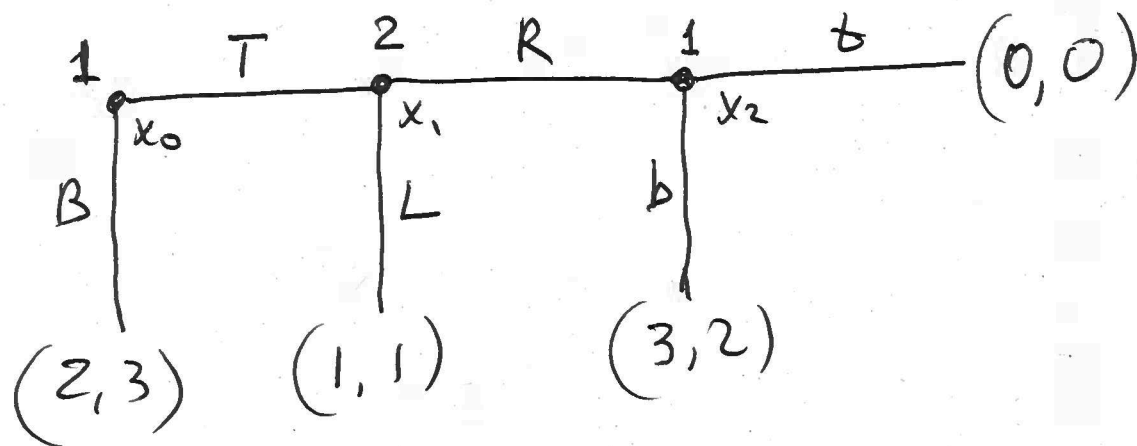
s_i is a complete set of instructions what to do at every node where player i needs to make a choice.

S_i = set of pure strategies of player i

$S = S_1 \times S_2 \times \dots \times S_N$ = set of pure strategy profiles of all players

$s = (s_1, s_2, \dots, s_N) \in S$ is a strategy profile

Example



$$i(x_0) = 1$$

$$i(x_1) = 2$$

$$i(x_2) = 1$$

$$A(x_0) = \{T, B\}$$

$$A(x_1) = \{L, R\}$$

$$A(x_2) = \{t, b\}$$

$$H_1 = \{x_0, x_2\}$$

$$H_2 = \{x_1\}$$

$$S_2 = \{L, R\}$$

$$S_1 = \{Tt, Tb, Bt, Bb\}$$

B is not a strategy!

Outcomes

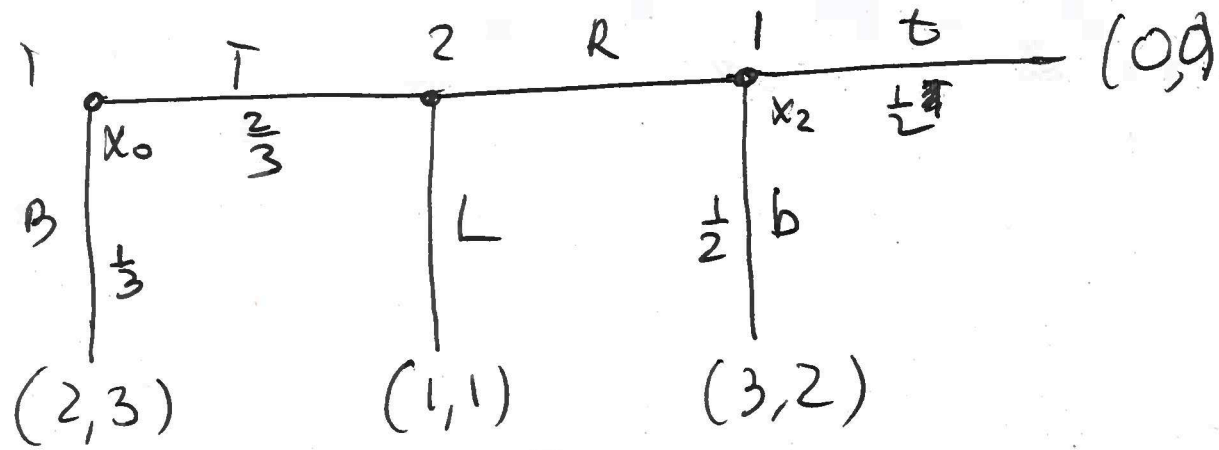
$o: S \rightarrow \Delta(Z)$ = the outcome of a strategy profile s .

$o(s)[z]$ = the probability that terminal node z is reached when strategy profile S is played.

Expected payoffs:

$$u_i: S \rightarrow \mathbb{R} \quad \text{where} \quad u_i(s) = \sum_{z \in Z} o(s)[z] \cdot u_i(z)$$

A mixed strategy $\sigma_i \in \Delta(S_i)$ is a probability distribution over pure strategies.



$$S_1 = \{ \cancel{Bb}, \cancel{Bt} \} \{ Tt, Tb, Bt, Bb \}$$

Example mixed str. σ_1 : $\sigma_1(Tt) = \sigma_1(Tb) = \sigma_1(Bb) = \frac{1}{3}$
 $\sigma_1(Bt) = 0$

$$\sigma_1(B) = \sigma_1(Bb) + \sigma_1(Bt) = \frac{1}{3} + 0 = \frac{1}{3}$$

$$\sigma_1(T) = \sigma_1(Tt) + \sigma_1(Tb) = \frac{1}{3} + \frac{1}{3} = \frac{2}{3}$$

$$\sigma_1(b|T) = \frac{\sigma_1(Tb)}{\sigma_1(T)} = \frac{1/3}{2/3} = \frac{1}{2}$$

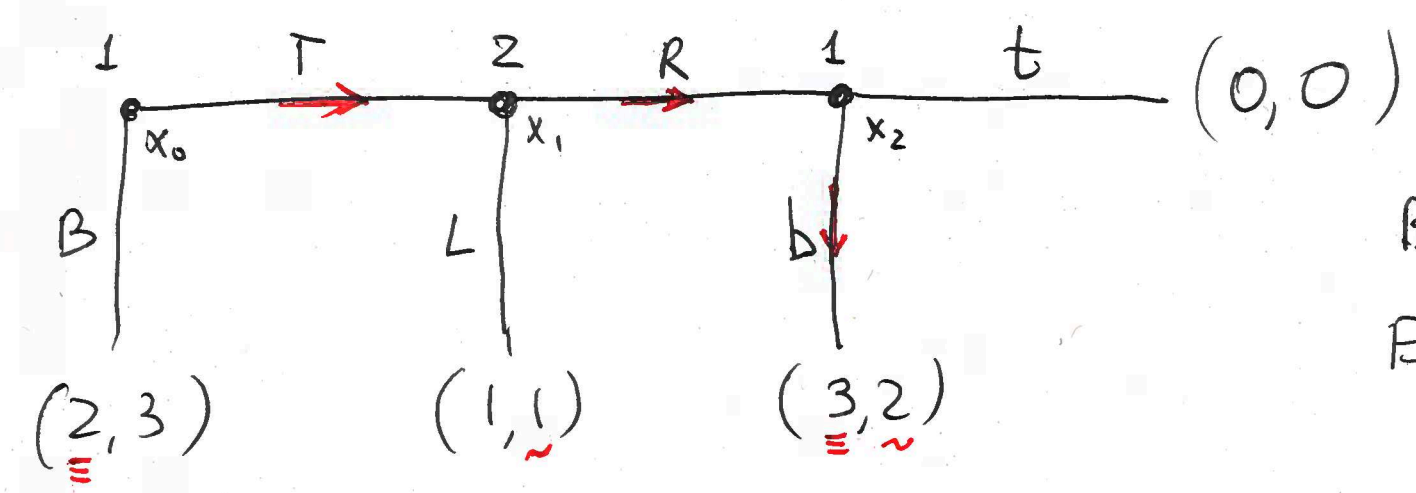
$$\sigma_1(t|T) = \frac{\sigma_1(Tt)}{\sigma_1(T)} = \frac{1/3}{2/3} = \frac{1}{2}$$

Bayes' Rule.

A behavior strategy $b_i \in \prod_{h \in H_i} \Delta(A(h))$ describes the probabilities of choosing actions in each of player i 's info. sets $h \in H_i$.
 $b_i(h)[a] =$ the probability of choosing action $a \in A(h)$ in the info. set h .

Proposition For games with perfect recall, restricting attention to behavior strategies is without loss of generality.

Backward Induction (BI)



Backward Induction
Equilibrium
(Tb, R)

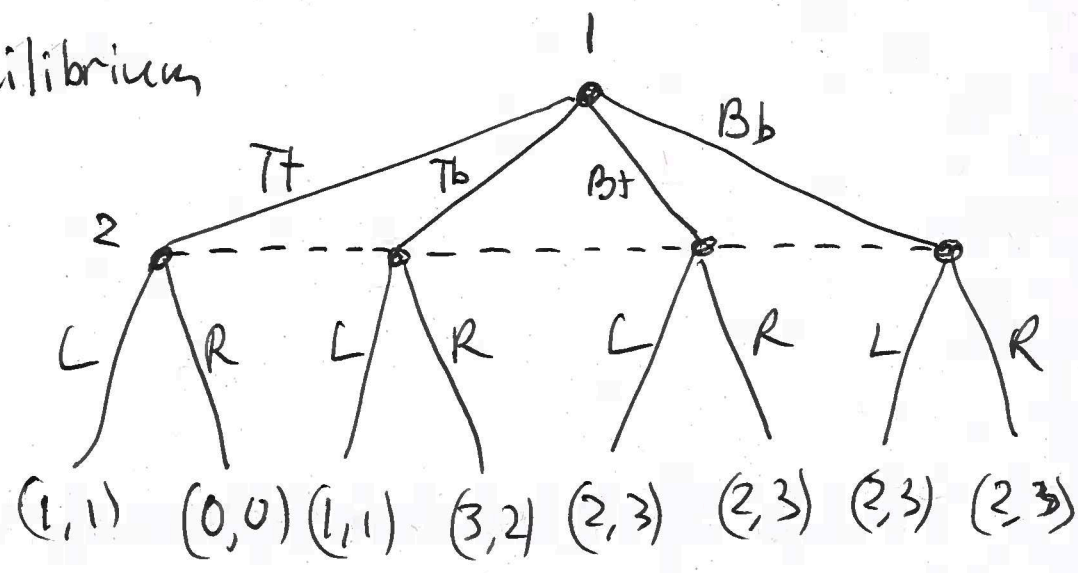
Game solvable by BI = Game of perfect information.

$$S_1 = \{Tt, Tb, Bt, Bb\}, \quad S_2 = \{L, R\}$$

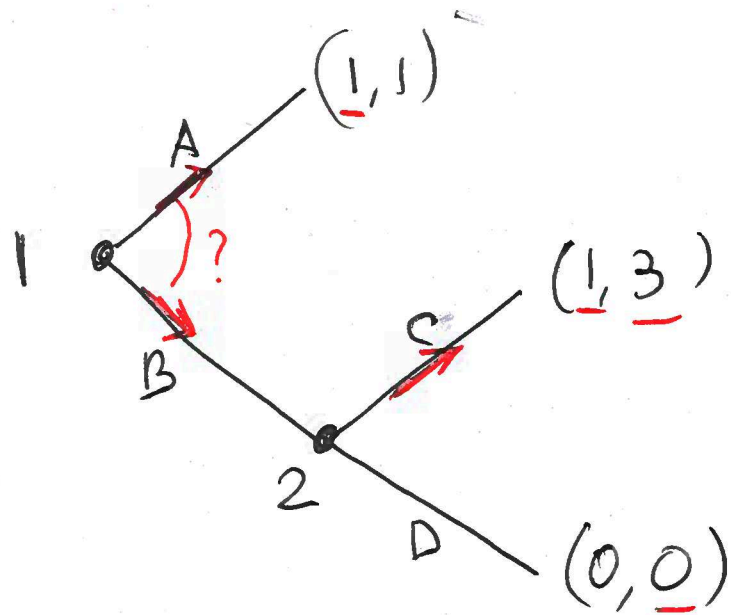
	L	R
Tt	1, 1*	0, 0
Tb	1, 1	*3, 2*
Bt	*2, 3*	2, 3*
Bb	*2, 3*	2, 3*

NE

BI Equilibrium



Proposition Every finite perfect information game has at least one BI equilibrium in pure strategies. Moreover, this BI eq. is unique, provided players have no ties in payoffs.



Pure strategy BI equilibria

(A, C) and (B, C)

The set of BI:

$$\{(\sigma_1(A), \sigma_1(B)), C\}$$

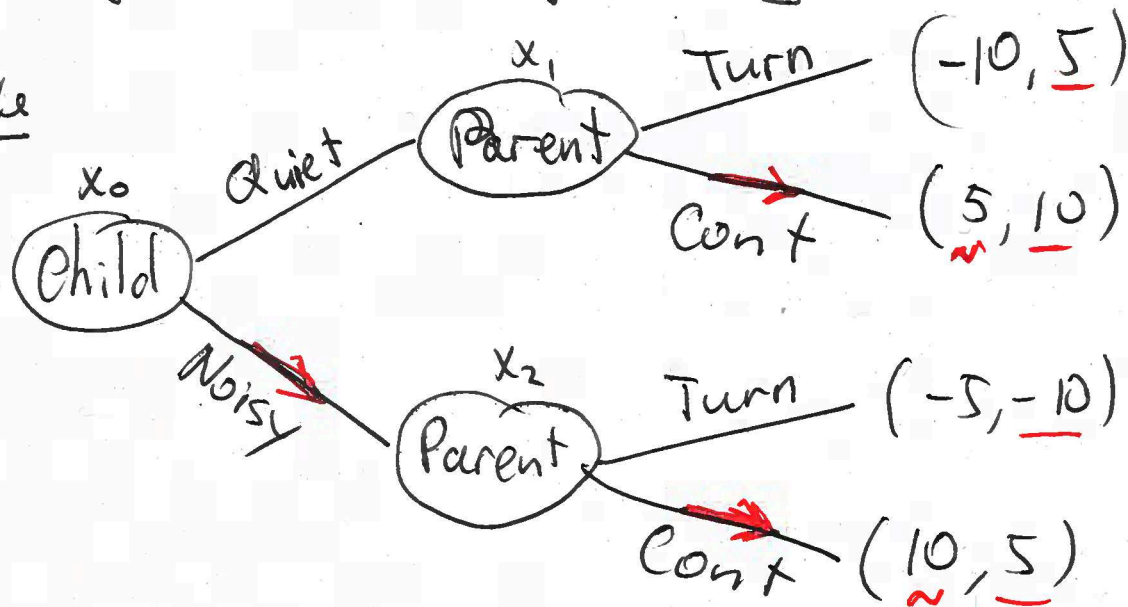
where $\sigma_1(A) + \sigma_1(B) = 1$

$$\sigma_1(A), \sigma_1(B) \geq 0.$$

Nash Equilibrium

A NE of an extensive-form game is a NE of the corresponding normal (matrix)-form game.

Example



$$S_1 = \{Q, N\}$$

$$S_2 = \{TT, TC, CT, CC\}$$

x_1, x_2

BI equilibrium:

$$(N, CC)$$

Non-credible threat
by the parent

	TT	TC	CT	CC
Q	-10, 5	-10, 5	* 5, 10 *	5, 10 *
N	* -5, -10	* 10, 5 *	-5, -10	* 10, 5 *

BI

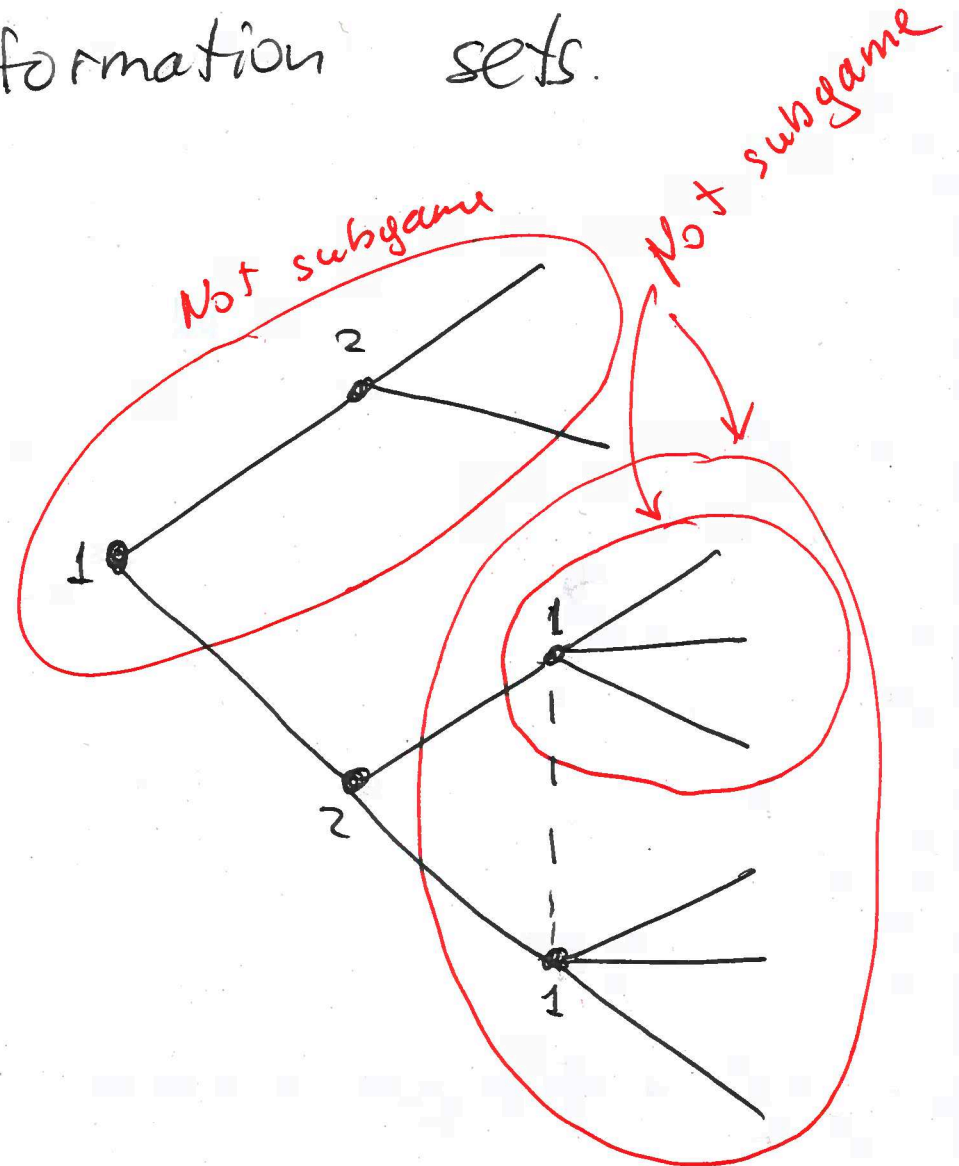
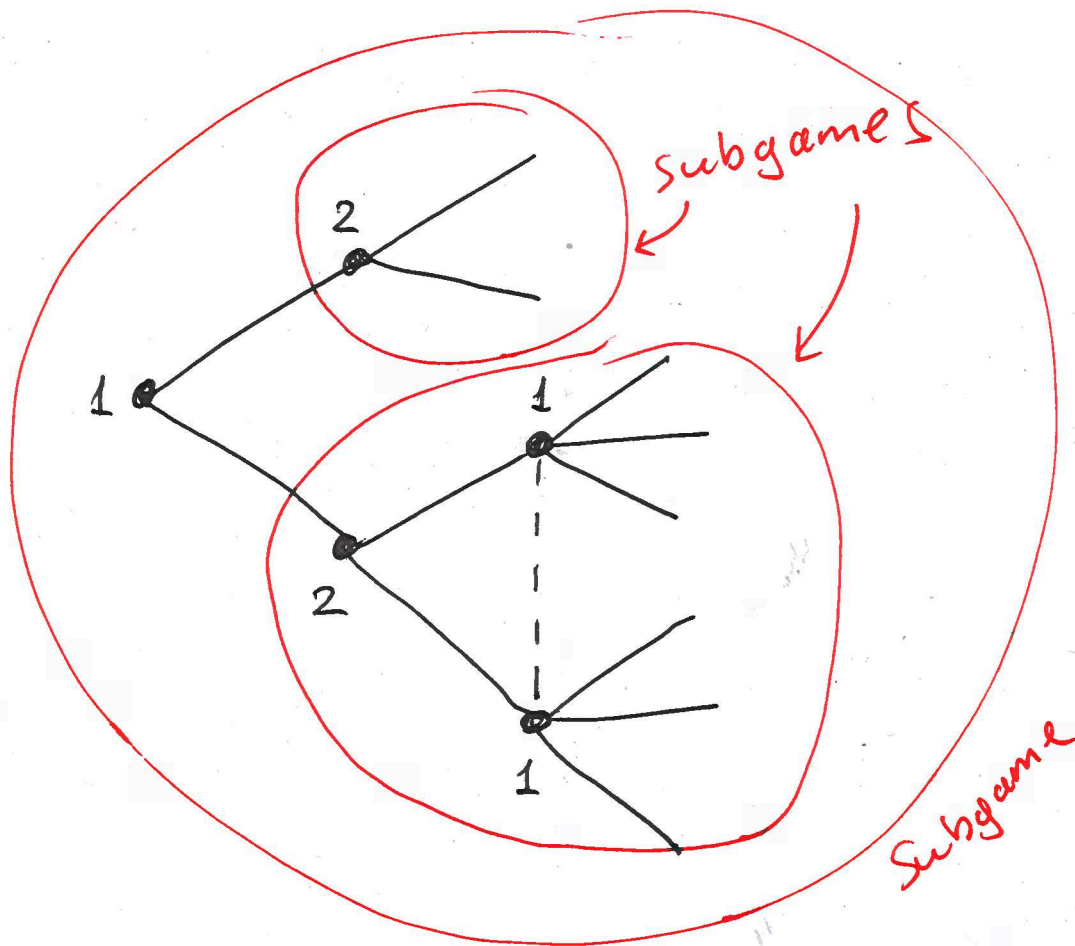
Proposition

Every BI equilibrium is a NE, but the converse need not be true.

Subgame Perfect Equilibrium

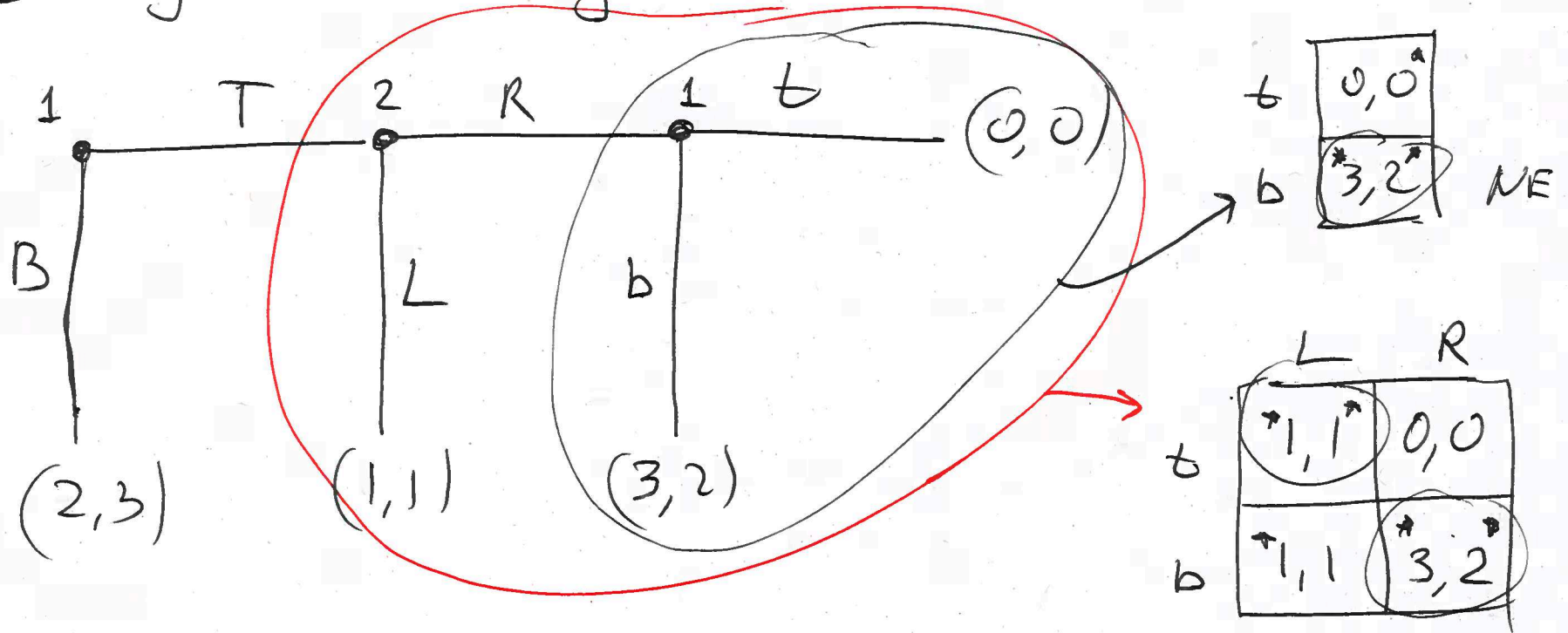
Def. A subgame is a subtree that

- (i) Starts at a decision node
- (ii) does not break information sets.

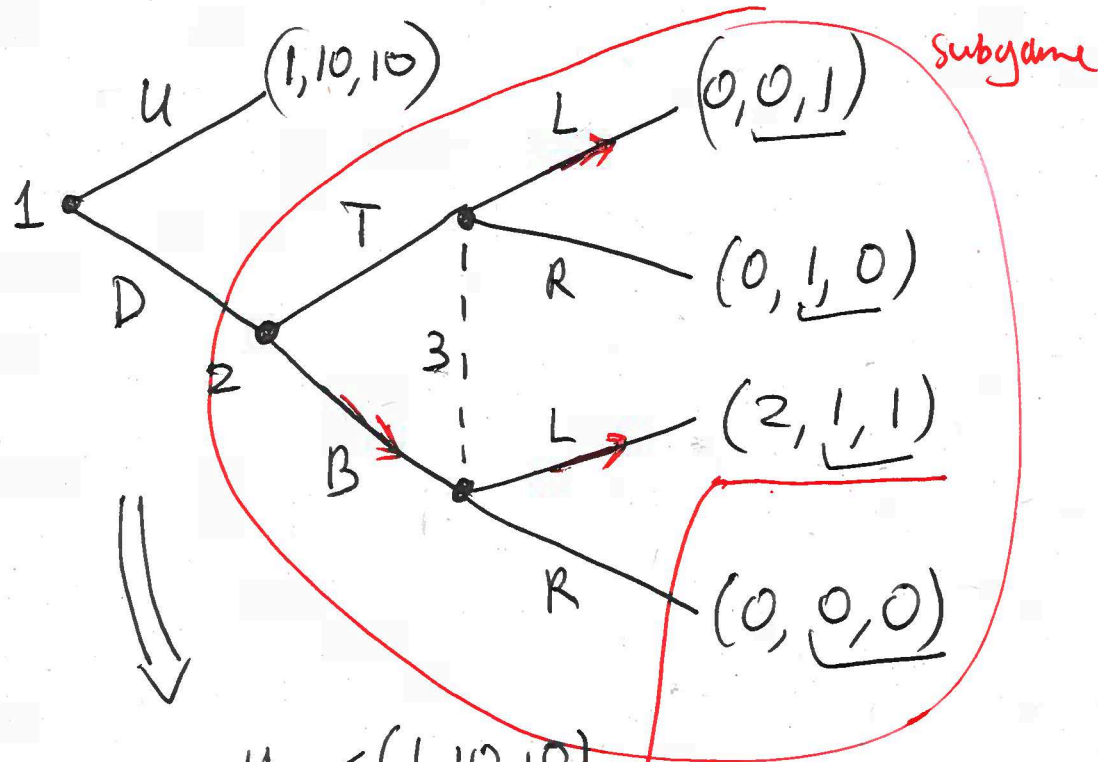


Def. A subgame perfect equilibrium (SPE)

is a profile of strategies such that their restriction to any subgame forms a NE of that subgame.



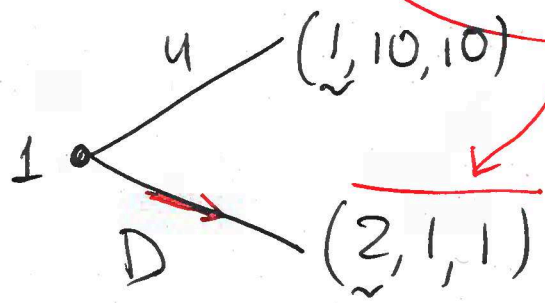
3-player example



Player 3

	L	R
Player 2 T	0, 1*	1, 0
Player 2 B	*1, 1*	0, 0

NE (B, L)
(ignoring player 1)



SPE: (D, B, L)

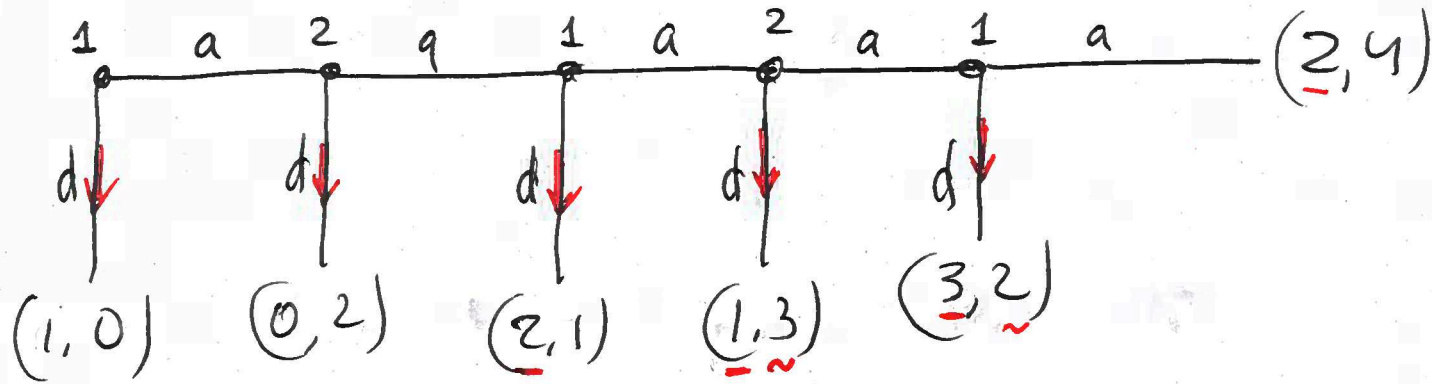
→ In games of perfect info, BI = SPE

→ In games that do not have any proper subgames, SPE = NE of the whole game

→ Every SPE is a NE of the whole game.

Problems with BI and SPE

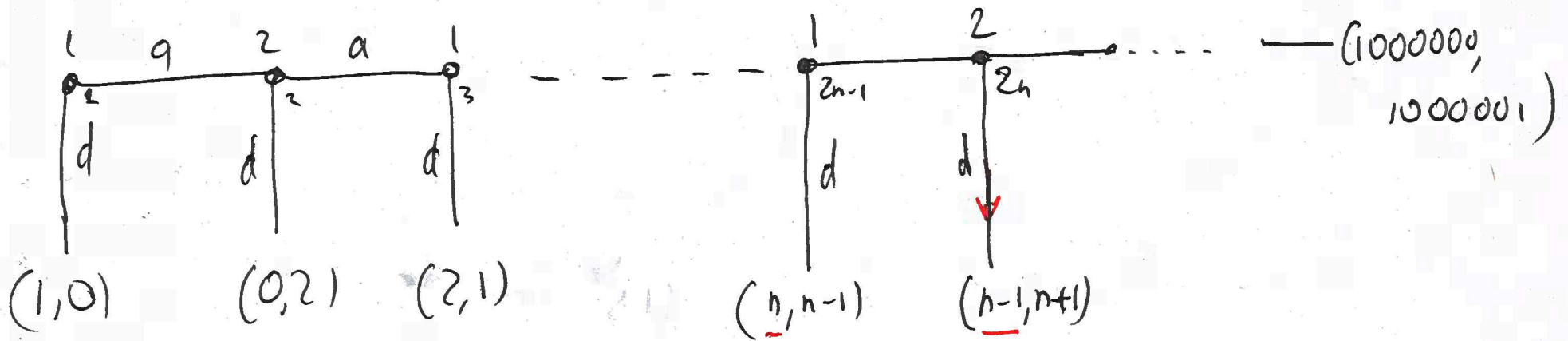
Centipede Game



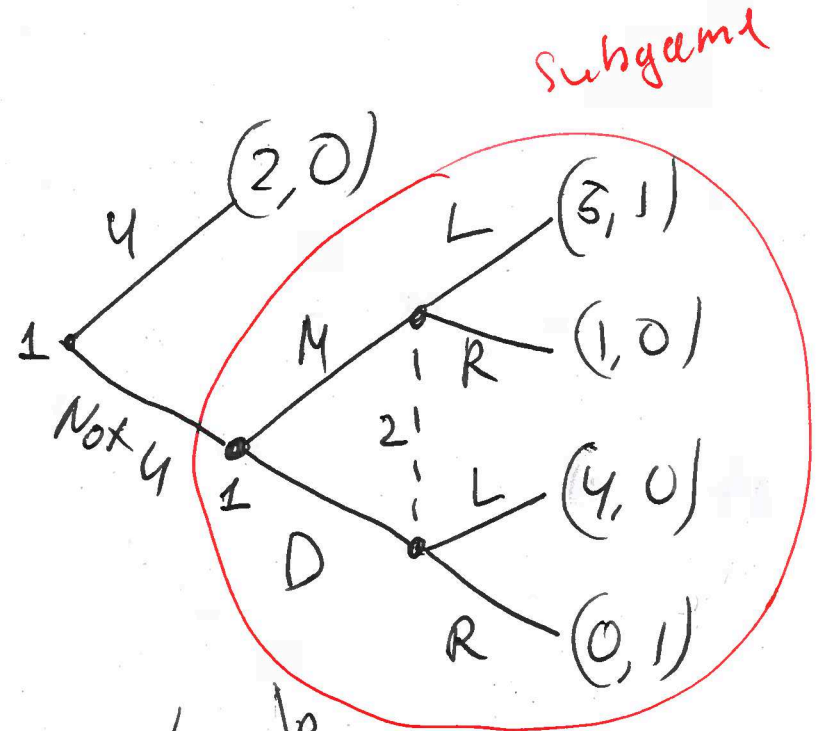
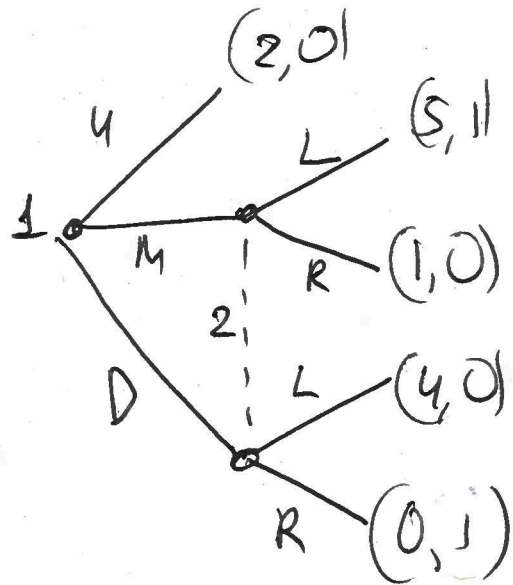
BI: (ddd, dd)

$$S_1 = \{aaa, aad, ada, \dots, dddd\}$$

$$S_2 = \{aa, ad, da, dd\}$$



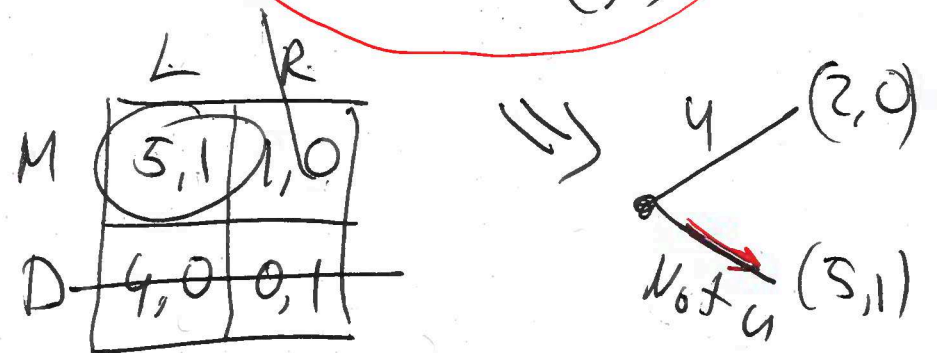
SPE: Might depend on seemingly irrelevant details



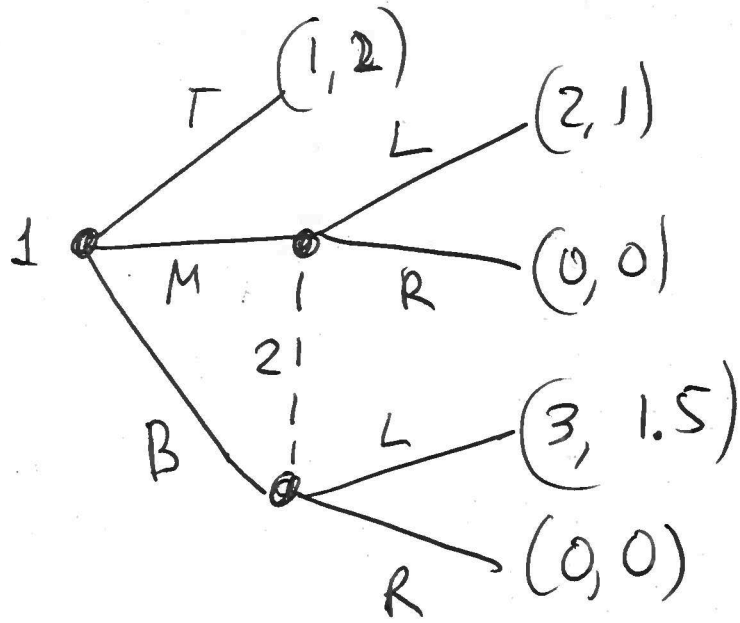
No subgames

	L	R
U	2,0*	*2,0*
M	*5,1*	1,0
D	4,0	0,1*

SPE



SPE: (Not U, M; L)
 → (5,1)

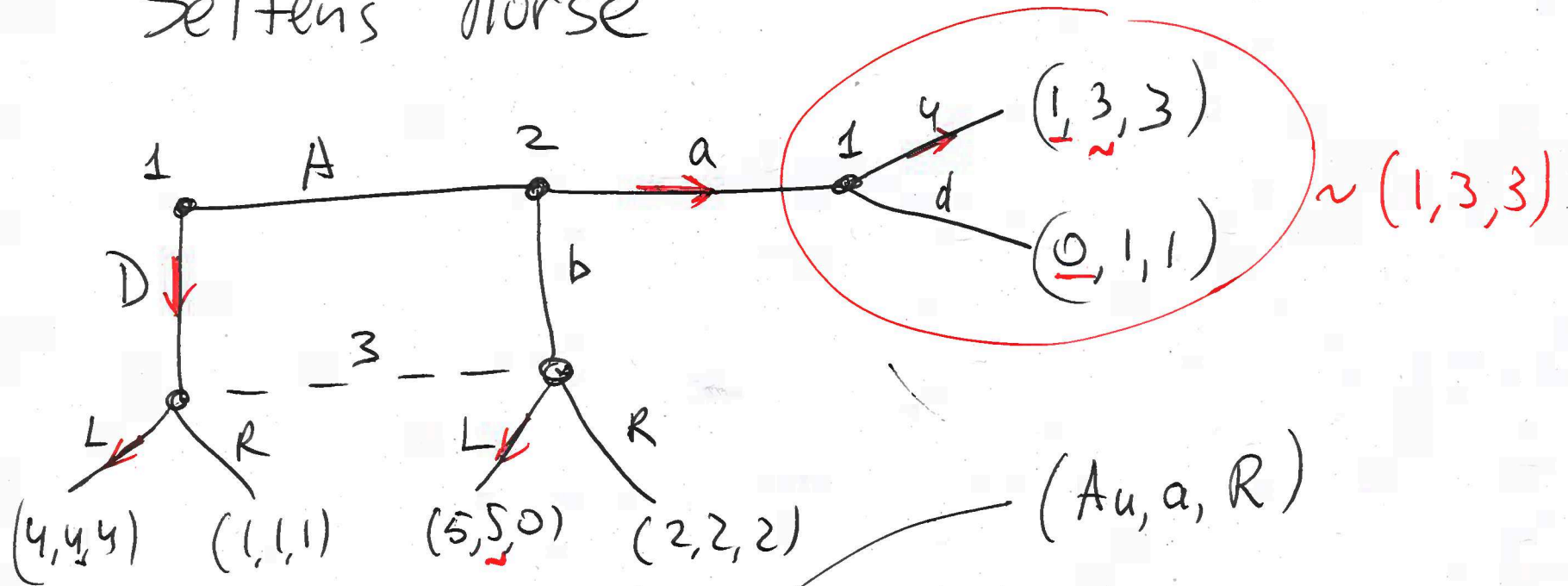


(T, R) is SPE

	L	R
T	1, 2*	*1, 2*
M	2, 1*	0, 0
B	*3, 1.5*	0, 0

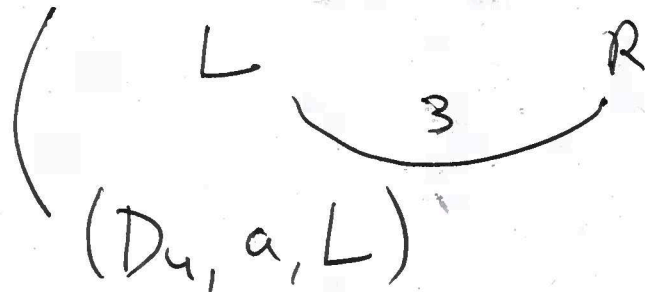
SPE

Selten's Horse



	a	b
A	1, 3, 3	5, 5, 0
D	4, 4, 4	4, 4, 4

	a	b
A	1, 3, 3	2, 2, 2
D	1, 1, 1	1, 1, 1



(A_u, a, R)