

Mixed Actions

= randomizations over pure actions

A mixed action ^{of player i} is a probability distribution

$$\sigma_i = (\sigma_i(a_1), \sigma_i(a_2), \dots, \sigma_i(a_k))$$

↑ probability to play action a_2

The support of a mixed action σ_i is

$$\text{Supp}(\sigma_i) = \{a_i \in A_i : \sigma_i(a_i) > 0\}$$

Expected utility from playing mixed action σ_i given beliefs p_j , $i, j \in \{1, 2\}$, is

$$U_i(\sigma_i, p_j) = \sum_{a_i \in A_i} \sigma_i(a_i) \cdot U_i(a_i, p_j) = \sum_{a_i \in A_i} \sum_{a_j \in A_j} \sigma_i(a_i) p_j(a_j) u_i(a_i, a_j)$$

belief p_j = mixed action p_j of player j

Consider player 1: $\max_{\sigma_1 \in \Delta(A_1)} u_1(\sigma_1, p_2)$

$$BR_1(p_2) = \operatorname{argmax}_{\sigma_1 \in \Delta(A_1)} u_1(\sigma_1, p_2)$$

Proposition A mixed action σ_1 is a BR to beliefs p_2 , so $\sigma_1 \in BR_1(p_2)$, if and only if every pure action in the support of σ_1 is a BR to p_2 . (In particular, every pure action in the support of σ_1 must have the same expected utility).

Interpretations of mixed action:

- literally randomize (to confuse the opponent)
- players have uncertainty (beliefs) about what others are doing.
- population of players with different pure actions

Def. Action $a_1 \in A_1$ is strictly dominated by mixed action ~~$\sigma_1 \in \Delta(A_1)$~~ $\sigma_1 \in \Delta(A_1)$ if

$$u(a_1, a_2) < u(\sigma_1, a_2) \quad \text{for all } a_2 \in A_2$$

Example

		L	R	
Player 1	T	3	0	$\frac{1}{2}$
	M	1	1	0
	B	0	3	$\frac{1}{2}$

} σ_2

M is strictly dominated by

$$\sigma_1 = \left(\frac{1}{2}, 0, \frac{1}{2} \right)$$

T M B

$$u_1(\sigma_1, L) = \frac{1}{2} \cdot 3 + 0 \cdot 1 + \frac{1}{2} \cdot 0 = 1.5 > u_1(M, L) = 1$$

$$u_1(\sigma_1, R) = \frac{1}{2} \cdot 0 + 0 \cdot 1 + \frac{1}{2} \cdot 3 = 1.5 > u_1(M, R) = 1$$

Proposition

In 2-player games, an action $a \in A_1$ is strictly dominated (possibly by a mixed action) if and only if a is NBR.

Iterative Deletion of Dominated Actions

Ex 1

		Rain God	
		Sun	Rain
Player 1	No umb.	(5, 3)	0, 0
	Umb.	1, 5	3, 2

ID(S) DA
 ↑ weakly
 strictly

- "Sun" strictly dominates "Rain"
→ delete "Rain"
- "No umb" strictly dominates "Umb"
→ delete "Umb"

Ex 2

		Pl. 2	
		L	R
Pl. 1	U	(8, 10)	4000, 9
	D	7, 6	6, 5

- L strictly dominates R
→ delete R
- U strictly dominates D
→ delete D

Common knowledge (of rationality, of payoffs)

An event is common knowledge (CK) if all players know that it is true, and all know that all know that it is true, and so on.

Example ID W DA
 \approx weakly

	D	E	F
A	1, 1	1, 1	2, 1
B	1, 1	0, 0	3, 1
C	1, 2	1, 3	1, 1

Case 1. Start with Pl. 1:

① A weakly dominates C \rightarrow delete C

② D weakly dominates E \rightarrow delete E

③ B ~~weakly dominates~~ A \rightarrow delete A
 \Rightarrow (B, D) and (B, F)

	D	F
B	1, 1	3, 1

Case 2. Start with Pl. 2:

① D weakly dominates F \rightarrow delete F

② A ~~weakly dominates~~ B \rightarrow delete B

③ E ~~weakly dominates~~ D \rightarrow delete D
 \Rightarrow (A, E) and (C, E)

	E
A	1, 1
C	1, 3

IDSDA \rightarrow order of deletion makes no difference

IDWDA \rightarrow order can matter.

A game is dominance solvable if the procedure of IDSDA (or the procedure of IDWDA regardless of the order of deletion) results in a unique outcome.

Rationalizability

An action is rationalizable if it cannot be (iteratively) deleted as NBR.
(the same as IDSDA in 2-player games)

Nash Equilibrium

An action profile $\sigma^* = (\sigma_1^*, \sigma_2^*, \dots, \sigma_n^*)$ is a Nash equilibrium if for each player i

$$u_i(\underbrace{\sigma_i^*}_{\sim}, \underbrace{\sigma_{-i}^*}_{\substack{\uparrow \\ \text{the other players}}}) \geq u_i(\underbrace{a_i}_{\sim}, \sigma_{-i}^*) \quad \text{for every } a_i \in A_i$$

That is, $\sigma_i^* \in BR(\sigma_{-i}^*)$ for each player i .

= "mutual best response".

		Rain	God	
		Sun	Rain	NE
PL. 1	No umb	*5, 3*	0, 0	
	Umb	1, 5*	*3, 2	

Proposition If $\sigma^* = (\sigma_1^*, \sigma_2^*)$ is a NE, then every action in its support } survives IDSDA
 } survives ID NBR A
 (Never-best-response)

	W	B	
W	$(2, 1)$	$(0, 0)$	q
B	$(0, 0)$	$(1, 2)$	$1-q$
	p	$1-p$	

NE
 (W, W)
 (B, B)

$$U_1(W, p) = 2 \cdot p + 0 \cdot (1-p) = 2p$$

$$U_1(B, p) = 0 \cdot p + 1 \cdot (1-p) = 1-p$$

For player 1 to mix between

W and B, the expected payoffs have to be the same:

$$2p = 1-p \Rightarrow p = \frac{1}{3}$$

$$U_2(W, q) = 1 \cdot q + 0 \cdot (1-q) = q$$

$$U_2(B, q) = 0 \cdot q + 2 \cdot (1-q) = 2 - 2q$$

$$q = 2 - 2q$$

$$\Rightarrow q = \frac{2}{3}$$

Mixed strategy NE:

$$\left(\begin{array}{c} q \quad 1-q \\ \left(\frac{2}{3}, \frac{1}{3} \right) \\ \underbrace{\quad \quad}_{\text{Player 1}} \\ W \quad B \end{array} \right), \left(\begin{array}{c} p \quad 1-p \\ \left(\frac{1}{3}, \frac{2}{3} \right) \\ \underbrace{\quad \quad}_{\text{Player 2}} \\ W \quad B \end{array} \right)$$

BR correspondence and NE

$$u_1(W, p) = 2p, \quad u_1(B, p) = 1 - p.$$

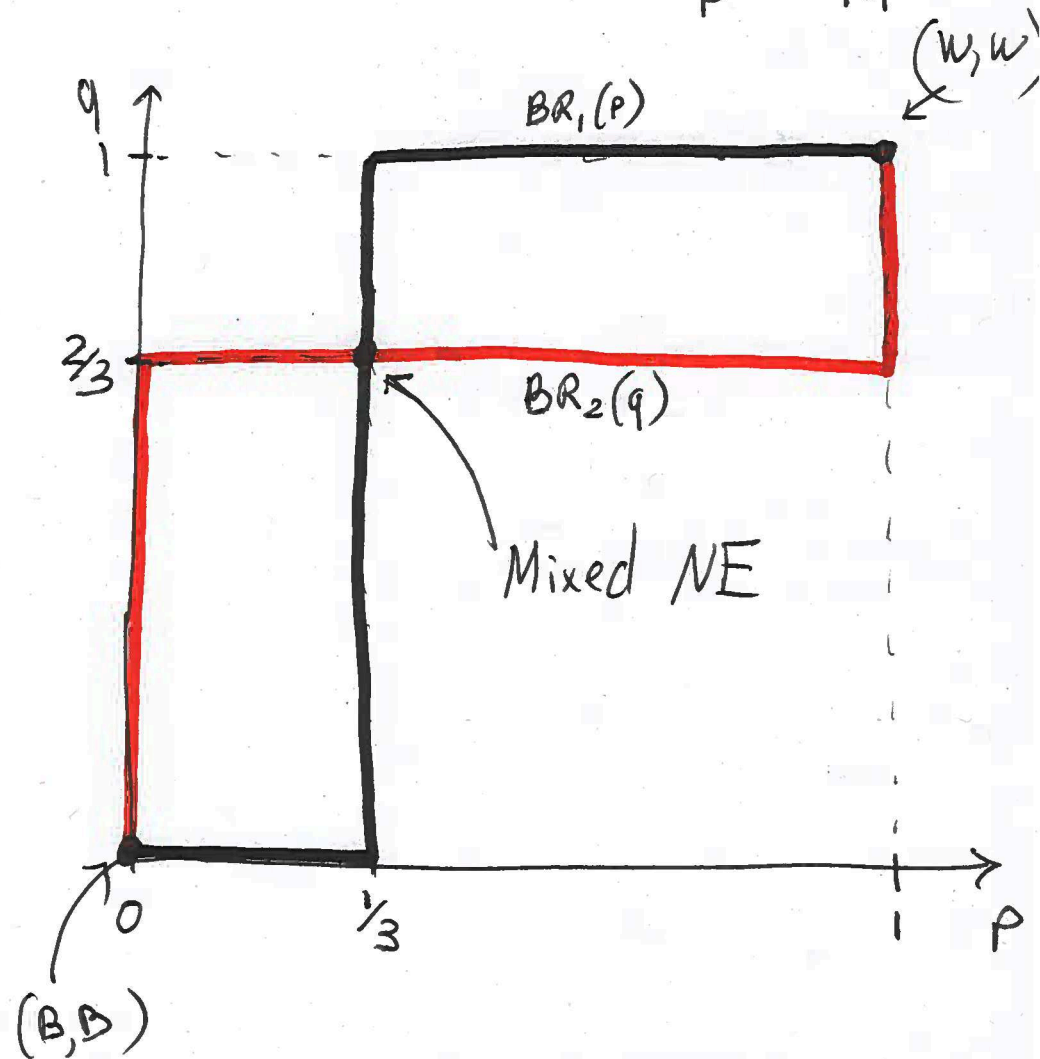
$$BR_1(p) = \begin{cases} W, & p > \frac{1}{3} \\ \text{anything}, & p = \frac{1}{3} \\ B, & p < \frac{1}{3} \end{cases}$$

$$BR_2(q) = \begin{cases} q = 1, & p > \frac{1}{3} \\ q \in [0, 1], & p = \frac{1}{3} \\ q = 0, & p < \frac{1}{3} \end{cases}$$

$$u_2(W, q) = q, \quad u_2(B, q) = 2 - 2q$$

$$BR_2(q) = \begin{cases} p = 1, & q > \frac{2}{3} \\ p \in [0, 1], & q = \frac{2}{3} \\ p = 0, & q < \frac{2}{3} \end{cases}$$

	W	B	
W	2, 1	0, 0	q
B	0, 0	1, 2	1-q
	p	1-p	



	L	R	
T	*10, 0	*5, 2*	q
B	*10, 11*	2, 11*	1-q
	p	1-p	

- ① T weakly dominates B } $\Rightarrow (T, R)$
 ② R — — — L }
 ① R — — — L } $\Rightarrow (T, R)$
 ② T — — — B }

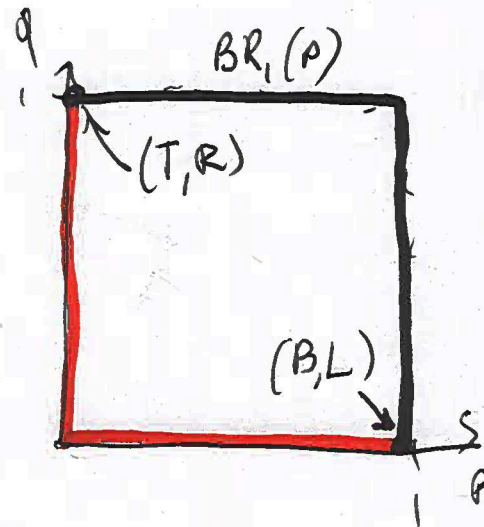
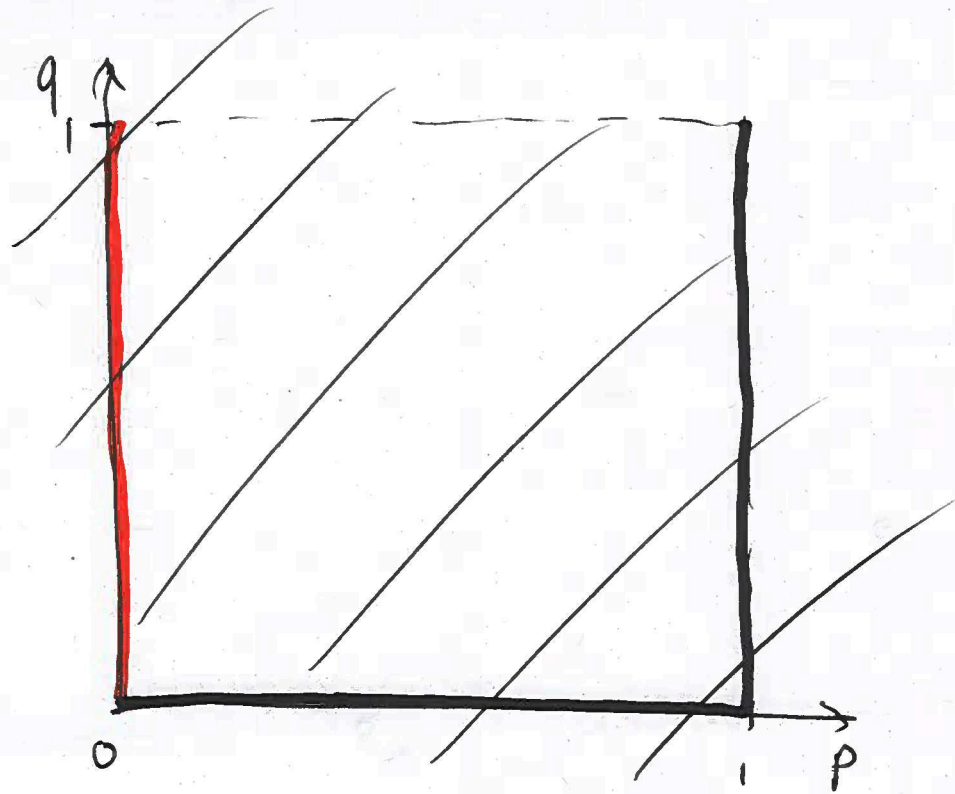
Dominance solvable!

$$\left. \begin{aligned} U_1(T, p) &= 10p + 5(1-p) \\ U_1(B, p) &= 10p + 2(1-p) \end{aligned} \right\}$$

$$\left. \begin{aligned} U_2(L, q) &= 0 \cdot q + 11(1-q) \\ U_2(R, q) &= 2 \cdot q + 11(1-q) \end{aligned} \right\}$$

$$BR_1(p) = \begin{cases} q=1, & p < 1 \\ q \in [0, 1], & p = 1 \end{cases}$$

$$BR_2(q) = \begin{cases} p=0, & q > 0 \\ p \in [0, 1], & q = 0 \end{cases}$$



Prisoner's Dilemma

	C	D
C	4, 4	0, 5
D	5, 0	1, 1

~~D~~ strictly dominates ~~C~~

Coordination game

	A	B
A	10, 10	0, 0
B	0, 0	1, 1

	L	R	
T	*3, 0	0, 3*	q
M	1, 1*	1, 1*	0
B	0, 3*	*3, 0	1-q
	p	1-p	

No pure NE.

$$U_1(T, p) = 3p + 0(1-p)$$

$$U_1(B, p) = 0 \cdot p + 3(1-p)$$

$$3p = 3(1-p)$$

$$p = \frac{1}{2}$$

$$U_2(L, q) = 0 \cdot q + 3(1-q)$$

$$U_2(R, q) = 3 \cdot q + 0(1-q)$$

$$3(1-q) = 3q,$$

$$q = \frac{1}{2}$$

NE in mixed actions

$$\left(\begin{array}{ccc} \frac{1}{2}, 0, \frac{1}{2} \\ T \quad M \quad B \end{array} \right), \left(\begin{array}{cc} \frac{1}{2}, \frac{1}{2} \\ L \quad R \end{array} \right)$$

Example 3 players

Player 1

		Player 2	
		X	Y
A	1, 1, 0	1, 0, 1	
B	1, 1, 1	0, 1, 1	

Player 2

		X	Y
A	1, 0, 1	1, 1, 0	
B	1, 1, 0	0, 1, 0	

NE in pure actions
(B, X, Box 1)

A weakly dominates B
delete B

NE \rightarrow (B, X, Box 1)
 \rightarrow (A, $(\frac{1}{2}, \frac{1}{2})$, $(\frac{1}{2}, \frac{1}{2})$)
x y Box 1 Box 2

Player 3

(A)

Player 3

		Box 1	Box 2	
Player 2	X	1, 0	0, 1	$\frac{1}{2}$
	Y	0, 1	1, 0	$\frac{1}{2}$
		$\frac{1}{2}$	$\frac{1}{2}$	

Divide the Dollar

$$A_1 = A_2 = \mathbb{R}_+$$

$$u_i(a_i, a_j) = \begin{cases} a_i, & a_i + a_j \leq 1 \\ 0, & a_i + a_j > 1 \end{cases}$$

$$a_j = x < 1 \Rightarrow BR_i(x) = \{1-x\}$$

$$a_j = x \geq 1 \Rightarrow BR_i(x) = \mathbb{R}_+$$

~~NE~~: any (a_i, a_j) :

$$a_i + a_j = 1 \quad \text{ex. } (0.4, 0.6)$$

any (a_i, a_j) :

$$\min\{a_i, a_j\} \geq 1$$

$$\text{ex. } (2, 1)$$

NE in weakly dominated actions.

